

Math 3032 (30 March 2021)

OH today 12-2pm.

HW 8 due today.

HW 9 (last one) due next week.

last week: gave
lots of criteria that
implied UFD.

e.g. PID, PID[x].

In \mathbb{Z} , $\mathbb{F}[x]$ we showed these were PIDs by
using long division. Essence was:

Given $a, b \neq 0$, can find q, r

s.t.

$$a = bq + r$$

and

either $r = 0$

or r smaller than b .]

don't need
unique.

in \mathbb{Z} , "smaller" could
mean $|r| < |b|$.

in $\mathbb{F}[x]$, it meant
 $\deg(r) < \deg(b)$.
[$\deg(0) = -\infty$]

Defn. Let R an integral domain.

A Euclidean norm on R is a function

$$v: R \setminus \{0\} \longrightarrow \mathbb{N} = \{0, 1, 2, \dots\}$$

s.t.

- if $a, b \in R \setminus \{0\}$ then $v(ab) \geq v(a)$
- $\forall a, b \in R \setminus \{0\}, \exists q, r$ s.t.

$$a = bq + r$$

and

$$\text{either } r = 0$$

$$\text{or } v(r) < v(b).$$

(R, v) is a
Euclidean domain

✓ Convince

Examples: $| - |$ is a Euclidean norm on \mathbb{Z} .

$\deg(-)$ is a Euclidean norm on $F[x]$.

Example

of a mult norm:

$| - |$ on \mathbb{Z} ,

$2^{\deg(-)}$ on $F[x]$.

$2^{-\infty} = 0$.

often but
not always,
 v will be multiplicative
 $\therefore v(ab) = v(a)v(b)$.
 $\because v(a) = 0$ then $a = 0$.

Let's review why Division \Rightarrow PID.

Thm: Euclidean domains are PIDs.

Pf: Suppose E a Euclidean domain, $I \subseteq E$ ideal.

Then if $I = \{0\}$, done. Otherwise, choose $b \in I$
s.t. $\nu(b) \leq \nu(r) \quad \forall r \in I - \{0\}$.

Then $\forall a \in I, \exists q, r$ s.t. $a = bq + r$.

Then $r \in I$. But if $r \neq 0$, then $\nu(r) < \nu(b)$
which is ruled out by assumption on b .

So $I = \langle b \rangle$.

\square .

Cor: Euclidean Domains are Noetherian UFDs.

The reason for the name "Euclidean" is because
in The Elements, Euclid gave an algorithm for
finding g.c.d.s. "Euclidean algorithm".

↑ fast/easy.

Reminder: In any UFD, any two elements a, b
have a gcd, well defined up to
association. $g = \gcd(a, b)$ if $g|a$ and $g|b$
and if $h|a$ and $h|b$ then $h|g$.

One way to find it: factor a, b into products
of primes, compare prime factors.

Factorization is slow/hard.

Euclid's algorithm: Let E be a Euclidean domain, $a, b \in E$

Set a_0 to be whichever of a, b has larger $v(-)$.

a_1 to be the other one.

$$\gcd(a, b) = \gcd(a_0, a_1)$$

associate

$$\exists q, r \text{ s.t. } a_0 = a_1 q + r$$

either $r=0$ or $v(r) < v(a_1)$.

Claim: $\gcd(a_0, a_1) \sim \gcd(a_1, r)$.

Pf: If $g \mid a_0$ and $g \mid a_1$ then

$$a_0 = m_0 g$$

$$a_1 = m_1 g$$

$$g \mid r \Rightarrow \gcd(a_0, a_1) \mid \gcd(a_1, r)$$

$$r = a_0 - a_1 q = (m_0 - m_1 q) g$$

Similarly if $g \mid a_1$ and $g \mid r$ then

$$g \mid a_0 \Rightarrow \gcd(a_1, r) \mid \gcd(a_0, a_1)$$

Set $a_2 := r$. Then divide a_2 into a_1 . Set $a_3 = \text{remainder}$.

$$\gcd(a, b) = \gcd(a_0, a_1) = \gcd(a_1, a_2) = \gcd(a_2, a_3) = \dots = \gcd(a_n, 0)$$

$$v(a_0) \geq v(a_1) > v(a_2) > \dots \geq 0 \quad \text{so alg. must stop} \quad \left. \vphantom{v(a_0)} \right\} a_n$$

Summary: $\gcd(a, b)$ is the last nonzero remainder after sequence of divisions.

Remark: Euclid's algorithm is special case of Gröbner's algorithm.

Note: Euclid's algorithm provides more than just a value of $\gcd(a, b)$.
"E-linear combination!"

It computes $m, n \in E$ s.t. $\gcd(a, b) = \overbrace{ma + nb}^{\text{set of elems}}$.

High powered reason: the computation was inside ideal $\langle a, b \rangle$.

$$a_0 = f_1 a_1 + a_2$$

$$a_1 = f_2 a_2 + a_3$$

...

$$a_{n-3} = f_{n-2} a_{n-2} + a_{n-1}$$

$$a_{n-2} = f_{n-1} a_{n-1} + a_n$$

$$(a_{n-1} = f_n a_n + 0.)$$

$$\Rightarrow \text{gcd} = a_n.$$

$$a_n = a_{n-2} - f_{n-1} (a_{n-3} - f_{n-2} a_{n-2})$$

$$= -f_{n-3} + \underbrace{(1 + f_{n-1} f_{n-2})}_{\text{orange}} a_{n-2} (a_{n-4} - \underbrace{f_{n-3} a_{n-3}}_{\text{orange}})$$

$$= (1 + f_{n-1} f_{n-2}) a_{n-4} + (-1)(1 + f_{n-1} f_{n-2})(-f_{n-3}) a_{n-3}.$$

$$a_0 - f_1 a_1 = a_2$$

⋮

$$a_{n-3} - f_{n-2} a_{n-2} = a_{n-1}$$

$$a_{n-2} - f_{n-1} a_{n-1} = a_n$$

If E is a Euclidean domain,

choice of v does not change $+$, \times .

Does constrain $+$, \times . (e.g. by forcing E to be a PID.)

v tells you about elements.

Prop: Let E be a Euclidean Domain. Then $u \in E$ is a unit iff $v(u) = v(1)$. And $v(1) \rightarrow$ the smallest value of v on $R \setminus \{0\}$.

Pf: If u is a unit, then for any $m \in E$,
 $v(m) = v(\underbrace{(mu^{-1})}_a \underbrace{u}_b) \geq v(b) = v(u)$.

If $v(u) = v(1)$, then write $1 = qu + r$.

Then $v(r) < v(u)$ impossible or $r=0$. So $1 = qu$.

Some examples of (Euclidean) domains from theory of numbers.

Defn: Let R be an integral domain. (E's any subring of \mathbb{C} .)

A multiplicative norm on R is a function $N: R \rightarrow \mathbb{Z}$ $\left(\begin{array}{c} \text{if want, can make positive} \\ \xrightarrow{1-1} \mathbb{N} \end{array} \right)$ by $|N(-)|$.

s.t. $N(ab) = N(a)N(b)$

$$\Rightarrow \begin{aligned} N(0^2) &= N(0) \\ N(1)^2 &= N(1^2) = N(1) \end{aligned}$$

and if $N(a) = 0$ then $a = 0$.

$$\text{so } N(0), N(1) \in \{0, 1\}.$$

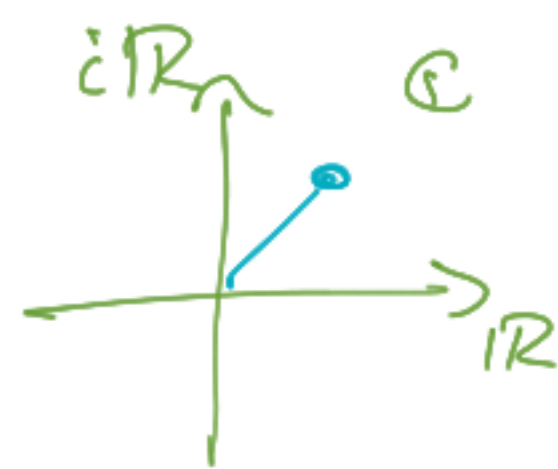
Remark: $|N(ab)| \geq |N(a)|$ if $a, b \neq 0$.

$$|N(a)| \cdot |N(b)| \geq 1$$

want demand
division.

Extremely nice example:

The Gaussian integers $\mathbb{Z}[i] \subseteq \mathbb{C}$



$$\mathbb{Z}^2 \stackrel{\text{as gp}}{\sim} \{a + bi \mid \text{s.t. } a, b \in \mathbb{Z}\}.$$

Set $N(a + bi) = |a + bi|^2 = a^2 + b^2.$

• integer? Yes because $a, b \in \mathbb{Z}.$

• $N(a + bi) = 0 \Rightarrow a + bi = 0$? Yes because

if $a^2 + b^2 = 0$, $a^2, b^2 \geq 0$ so must have

$$a^2 = b^2 = 0. \quad \text{So } a = b = 0.$$

• multiplicative? Set $\overline{a + ib} = a - ib.$

Given $\alpha \in \mathbb{C}$ set $\bar{\alpha} := a - ib$.

"
 $a + ib$

if $\alpha \in \mathbb{Z}[i]$ then so is $\bar{\alpha}$.

$\alpha \mapsto \bar{\alpha}$ is an involution of $\mathbb{Z}[i]$.
 $\bar{\bar{\alpha}} = \alpha$

and a ring hom.

$$\overline{\alpha + \beta} = \bar{\alpha} + \bar{\beta} \quad \overline{\alpha\beta} = \bar{\alpha}\bar{\beta}.$$

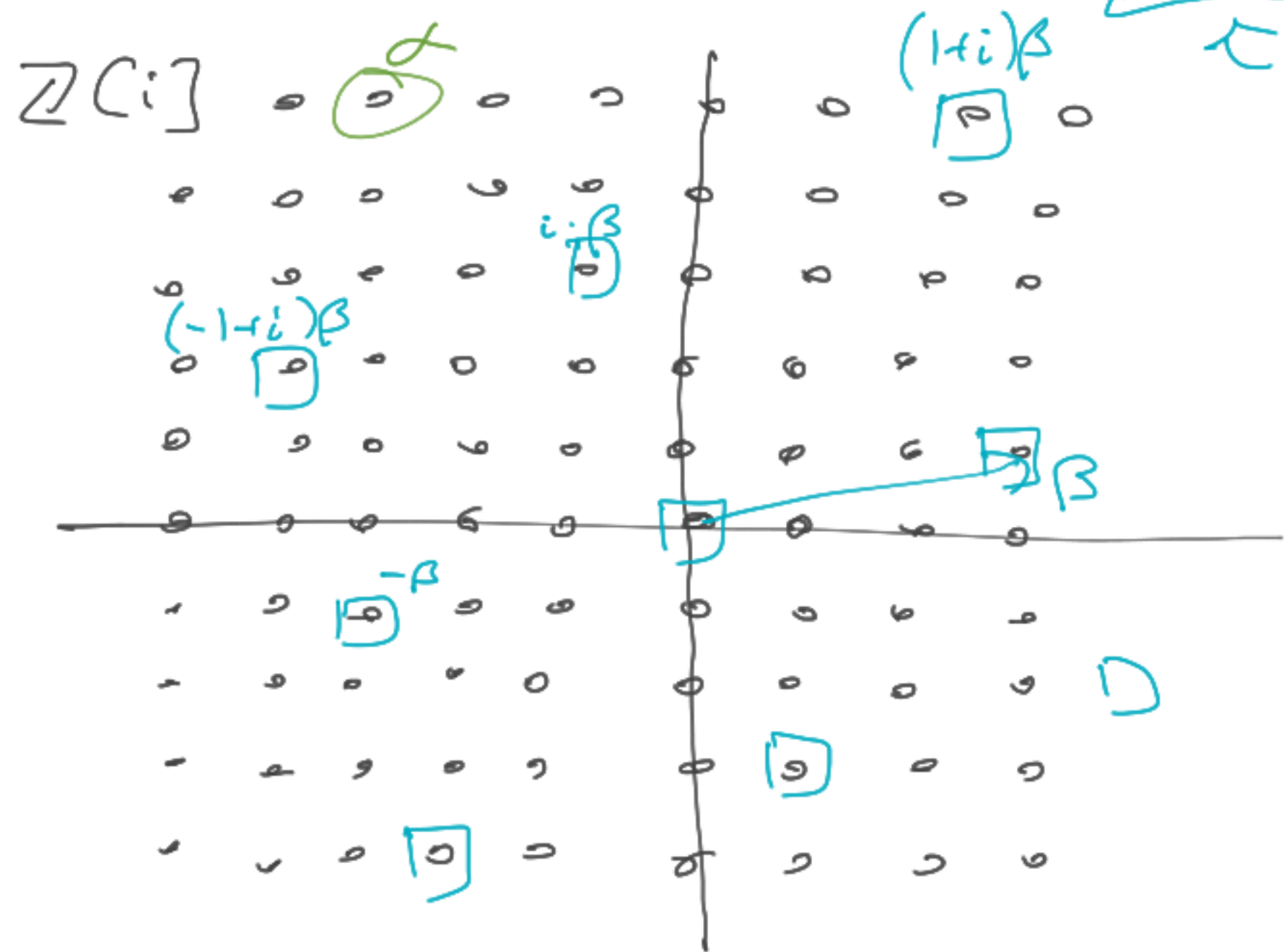
Recognize: $N(\alpha) = \alpha \cdot \bar{\alpha}$.

$$\begin{aligned} \text{Then } N(\alpha\beta) &= \alpha\beta \overline{\alpha\beta} = \alpha\beta \bar{\alpha}\bar{\beta} \\ &= \alpha \bar{\alpha} \beta \bar{\beta} = N(\alpha)N(\beta). \end{aligned}$$

com.

Thm: $(\mathbb{Z}[i], N)$ is Euclidean.

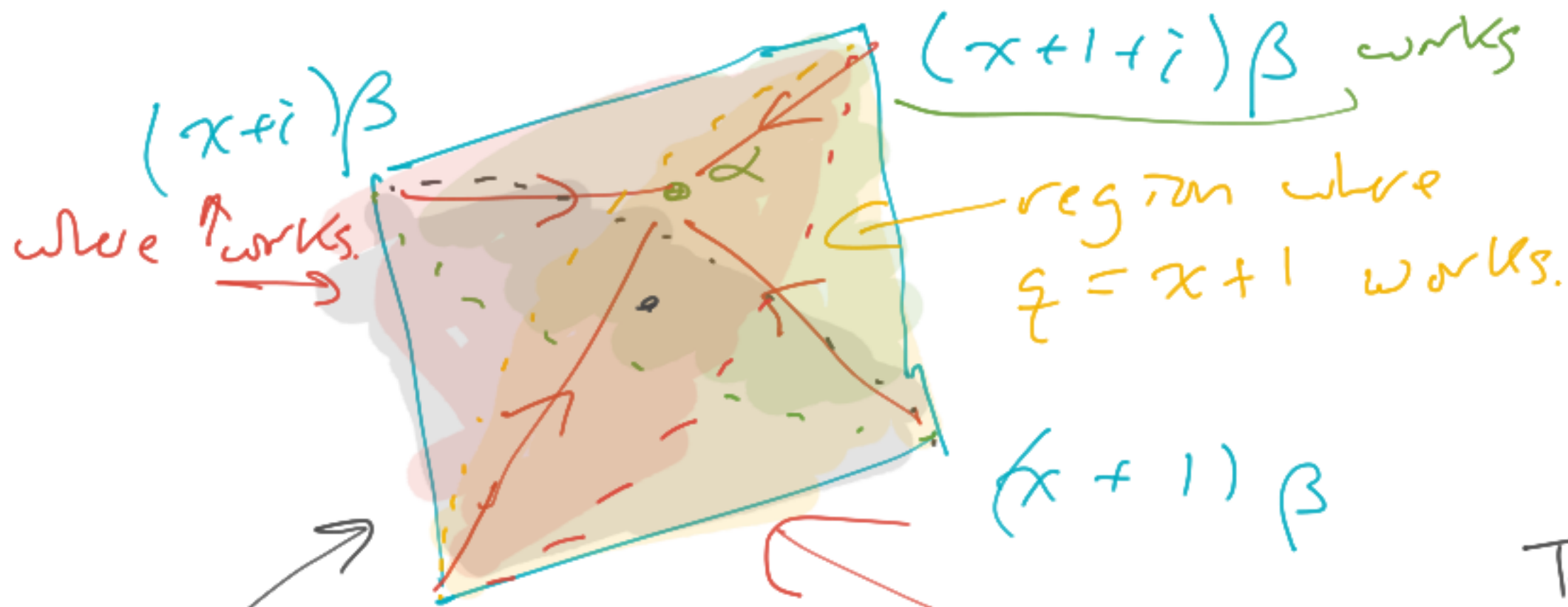
pf: We need to verify: $\forall \alpha, \beta \in \mathbb{Z}[i], \beta \neq 0$
 $\exists \gamma, r$ s.t. $\alpha = \gamma\beta + r$ and $N(r) < N(\beta)$.



as γ varies, these
comprise $\langle \beta \rangle$.

$\langle \beta \rangle$ is a
rescaled rotated
square grid.





So some square contains α (perhaps on boundary).

Try to set

$$\eta \in \{x, x+1, x+i, x+i\}$$

r will be one of these vectors.

We win if at least one of them has $|r|^2 \neq |\beta|^2$.

all the α 's for which $|\alpha - x\beta|^2 \leq |\beta|^2$.

i.e. ones where $\eta = x$ works.

Punchline: These quarter circles cover the square. \square

Cor: $\mathbb{Z}[i]$ is a UFD.

But factorization in $\mathbb{Z}[i]$
 \neq factorization in \mathbb{Z} .

$$\begin{aligned} 5 &= (2+i)(2-i) \\ &= (1+2i)(1-2i) \end{aligned}$$

so 5 is not prime.

$$i \cdot (-i) = 1.$$

$$2+i = i \cdot (1-2i) \quad i \text{ is unit.}$$