

See arXiv:1502.06526, joint w/ C. Scheimbauer.

1. Functorial qfts: Schrödinger v. Heisenberg

Spacetime categories: QFTs associate algebraic data (“partition functions”, “OPEs”, ...) to geometric data (“spacetimes” equipped with some geometric structure). This should be *local* when spacetimes are cut and pasted.

How to organize “locality”? One way is to build an (∞, n) -category SPACETIMES whose k -dim morphisms are k -manifolds surrounded by germs of n -dimensional manifolds equipped with desired geometric structure.

Henceforth: fix choice of SPACETIMES.

Schrödinger-picture qft: The *Schrödinger picture* of quantum mechanics uses systems described by Hilbert spaces and linear maps between them. Atiyah (et al.): a *Schrödinger picture n -dimensional qft* is a symmetric monoidal (∞, n) -functor $Z : \text{SPACETIMES} \rightarrow \text{VECT}_n$, where $\text{VECT}_n = \{\text{linear } (\infty, n-1)\text{-categories}\}$ is an (∞, n) -category whose $(n-1)$ -morphisms are vector spaces (maybe topological, chain complexes, etc.).

Problem: Hilbert spaces aren’t physical. “Space of states” is $\mathbb{P}\mathcal{H}$, not $\mathcal{H} \in \text{VECT}$. But do want \otimes, \oplus .

Solution: Pointed category $(\mathcal{H}, \text{VECT})$ has

$$\text{Aut}((\mathcal{H}, \text{VECT})) \cong \text{PGL}(\mathcal{H}) = \text{Aut}(\mathbb{P}\mathcal{H}).$$

Remark: Often, $(\mathcal{H}, \text{VECT}) \simeq (A, \text{MOD}_A)$ in bicategory of pointed categories, i.e. is *affine*.

Heisenberg-picture qft: Should assign pointed categories to spaces. What about spacetimes, etc.?

Heisenberg = relative: A pointed category (C, \mathcal{C}) is a morphism $C : \mathbb{1} \rightarrow \mathcal{C}$, where $\mathbb{1}$ is the monoidal unit in the bicategory of categories. So *Heisenberg-picture qft* could be a “morphism” $Z : \mathbb{1} \Rightarrow \mathcal{Z}$ where $\mathbb{1} : \text{SPACETIMES} \rightarrow \text{VECT}_{n+1}$ is constant functor and $\mathcal{Z} : \text{SPACETIMES} \rightarrow \text{VECT}_{n+1}$ is a “categorified Schrödinger picture qft,” called the *twist*.

Guess: “Morphism” = “sym \otimes natural transformation.”

Exercise: For topological 1-bordisms cat BORD_1 , all sym \otimes “strong” nat trans between functors $\text{BORD}_1 \rightarrow \mathcal{C}$ are isomorphisms, for any sym \otimes bicategory \mathcal{C} .

Defn: An *(op)lax transformation* between functors of bicategories is like a nat trans, but “naturality” is just a morphism, not nec. invertible. In a *strong* transformation, naturality is imposed by isomorphisms.

2. What is an (∞, n) -category?

\exists many equivalent models of “ (∞, n) -category.” I will describe “complete n -fold Segal space” model.

Warm-up: A *category* is a simplicial set $\mathcal{C}_\bullet : \Delta^{\text{op}} \rightarrow \text{SETS}$ such that (“Segal condition”) $\forall k \geq 1$, the map

$$\mathcal{C}_k \rightarrow \underbrace{\mathcal{C}_1 \times_{\mathcal{C}_0} \mathcal{C}_1 \times_{\mathcal{C}_0} \cdots \times_{\mathcal{C}_0} \mathcal{C}_1}_{k \text{ times}}$$

is an iso. (Use $\{(i-1, i) : \Delta^1 \rightarrow \Delta^k\}_{i=1}^k$.)

A *functor* is a map of simplicial sets. BUT: this gets WRONG notion of “equivalence” (no isos of functors).

Defn: An $(\infty, 1)$ -category is a simplicial topological space $\mathcal{C}_\bullet : \Delta^{\text{op}} \rightarrow \text{SPACES}$ such that $\forall k \geq 1$, the map

$$\mathcal{C}_k \rightarrow \underbrace{\mathcal{C}_1 \times_{\mathcal{C}_0}^h \mathcal{C}_1 \times_{\mathcal{C}_0}^h \cdots \times_{\mathcal{C}_0}^h \mathcal{C}_1}_{k \text{ times}}$$

is a homotopy equivalence ($\times_{\mathcal{C}_0}^h$ is homotopy fibered product), AND MOREOVER (“completeness”) every (family of) invertible-up-to-homotopy 1-morphism(s) is homotopic to an (family of) identity 1-morphism(s).

A *strict functor* is a map of simplicial spaces. Note: $\text{hom}(\mathcal{B}_\bullet, \mathcal{C}_\bullet)$ is automatically a space. A functor is an *equivalence* if it is levelwise so. Localize at equivalences. This gets the RIGHT notion.

Remark: Model category theory makes this all clean. Rather than localizing, work just with $(\infty, 1)$ -categories that as simplicial spaces are cofibrant for projective model structure. Or work just with fibrants for injective model.

Defn: An (∞, n) -category is an n -fold simplicial space $\mathcal{C}_\bullet : (\Delta^{\times n})^{\text{op}} \rightarrow \text{SPACES}$ which is $(\infty, 1)$ -cat in each direction (i.e. an $(\infty, 1)$ -cat internal to ... internal to $(\infty, 1)$ -cats), AND MOREOVER $\forall i \leq n, \forall \vec{k} \in \mathbb{N}^i, \mathcal{C}_{\vec{k}, 0, \dots, 0} \cong \mathcal{C}_{\vec{k}, 0, 0, \dots, 0}$ (i.e. at every level, categories of objects are actually spaces). So \mathcal{C} is built from the spaces $\mathcal{C}_{(i)} = \underbrace{\mathcal{C}_{1, \dots, 1, 0, \dots, 0}}_{i \text{ 1s}} = \{i\text{-morphisms in } \mathcal{C}\}$.

Strict functors are maps of n -fold simplicial spaces. *Equivalences* are levelwise. Localize at those. Or use model category theory.

Defn: A *symmetric monoidal (∞, n) -category* is $\mathcal{C}_\bullet^\otimes(\bullet) : (\Delta^{\times n} \times \Gamma)^{\text{op}} \rightarrow \text{SPACES}$ (Segal’s Γ is opposite to {finite pointed sets}), which is (∞, n) - in the $\Delta^{\times n}$ direction, and $\forall k, \mathcal{C}_\bullet^\otimes(k) \rightarrow \mathcal{C}_\bullet^\otimes(1)^{\times k}$ is a homotopy equivalence.

Strict functors, equivalences, localize, model categories.

3. Walking morphisms and the main definition

For (∞, n) -categories \mathcal{B}, \mathcal{C} , let $\text{maps}(\mathcal{B}, \mathcal{C}) = \text{derived space of functors} = \text{mapping space in the homotopy category} = \text{hom}(\text{cofibrant replacement of } \mathcal{B}, \text{ fibrant replacement of } \mathcal{C})$.

$\{(\infty, n)\text{-categories}\}$ is *Cartesian closed*: $\forall \mathcal{B}, \mathcal{C}, \exists$ an (∞, n) -category $\underline{\text{maps}}(\mathcal{B}, \mathcal{C})$ s.t.

$$\text{maps}(-, \underline{\text{maps}}(\mathcal{B}, \mathcal{C})) \simeq \text{maps}((-) \times \mathcal{B}, \mathcal{C}).$$

Its space of objects $\underline{\text{maps}}(\mathcal{B}, \mathcal{C})_{(0)}$ is just $\text{maps}(\mathcal{B}, \mathcal{C})$. Its space of 1-morphisms $\underline{\text{maps}}(\mathcal{B}, \mathcal{C})_{(1)}$ is the space of *strong natural transformations*.

Defn: The *walking 0-morphism* is $\Theta^{(0)} = \{\bullet\}$. The *walking 1-morphism* is $\Theta^{(1)} = \{\bullet \rightarrow \bullet\}$. The *walking 2-morphism* is $\Theta^{(2)} = \left\{ \bullet \begin{array}{c} \Downarrow \\ \bullet \end{array} \bullet \right\}$. Etc. These satisfy:

$$\mathcal{C}_{(i)} \simeq \text{maps}(\Theta^{(i)}, \mathcal{C}).$$

$$\begin{aligned} \text{So } \underline{\text{maps}}(\mathcal{B}, \mathcal{C})_{(1)} &\simeq \text{maps}(\Theta^{(1)}, \underline{\text{maps}}(\mathcal{B}, \mathcal{C})) \\ &\simeq \text{maps}(\Theta^{(1)} \times \mathcal{B}, \mathcal{C}) \simeq \text{maps}(\mathcal{B}, \underline{\text{maps}}(\Theta^{(1)}, \mathcal{C})). \end{aligned}$$

I.e. a strong transformation η assigns:

$$\begin{aligned} B \in \mathcal{B}_{(0)} &\mapsto \eta(B) \in \underline{\text{maps}}(\Theta^{(1)}, \mathcal{C})_{(0)} \simeq \mathcal{C}_{(1)} \\ b \in \mathcal{B}_{(1)} &\mapsto \eta(b) \in \underline{\text{maps}}(\Theta^{(1)}, \mathcal{C})_{(1)} \\ &\simeq \text{maps}(\Theta^{(1)} \times \Theta^{(1)}, \mathcal{C}) \end{aligned}$$

This gives *strong* transformations because

$$\Theta^{(1)} \times \Theta^{(1)} = \begin{array}{ccc} \bullet & \xrightarrow{\quad} & \bullet \\ \downarrow & \nearrow & \downarrow \\ \bullet & \xrightarrow{\quad} & \bullet \end{array}.$$

Strategy: Replace $\text{maps}(\Theta^{(1)}, \mathcal{C})$ by an (∞, n) -category with the same objects (=1-morphisms in \mathcal{C}), but with 1-morphisms = diagrams of shape $\begin{array}{ccc} \bullet & \xrightarrow{\quad} & \bullet \\ \downarrow & \nearrow & \downarrow \\ \bullet & \xrightarrow{\quad} & \bullet \end{array}$ in \mathcal{C} .

This shape is not symmetric. We want it to be a 1-morphism between arrows in \mathcal{C} , so that a “transformation” assigns an arrow to each object and one of these squares to each morphism. So let’s decide: in “ \mathcal{C}^\rightarrow ,” the objects are “horizontal” arrows, and a 1-morphism is a box reading down, with source and target (of the 1-morphism) the “vertical” source and target s_v, t_v . In “ \mathcal{C}^\downarrow ,” the objects are “vertical” arrows; the source and target of a box-as-1-morphism are the “horizontal” source and target s_h, t_h .

Note that then there are two maps $s_v, t_v : \mathcal{C}^\downarrow \rightrightarrows \mathcal{C}$ that take any vertical arrow or any diagram just to its top or bottom parts. Similarly $s_h, t_h : \mathcal{C}^\rightarrow \rightrightarrows \mathcal{C}$.

Defn: Suppose we have defined such \mathcal{C}^\rightarrow and \mathcal{C}^\downarrow . A *lax transformation* $\eta : F \rightrightarrows G$ between $F, G : \mathcal{B} \rightrightarrows \mathcal{C}$ is a functor $\eta : \mathcal{B} \rightarrow \mathcal{C}^\downarrow$ with $s_v \circ \eta = F$ and $t_v \circ \eta = G$. An *oplax transformation* is $\eta : \mathcal{B} \rightarrow \mathcal{C}^\rightarrow$ (c.f. §6).

In fact, \mathcal{C}^\rightarrow and \mathcal{C}^\downarrow are the horizontal and vertical 1-morphism categories of a *double* (∞, n) -category \mathcal{C}^\square , i.e. $\mathcal{C}^{\rightarrow\downarrow} = \mathcal{C}^{\square}_{(1);\bullet}$ and $\mathcal{C}^{\downarrow\rightarrow} = \mathcal{C}^{\square}_{\bullet;(1)}$. Basic data of an (∞, n) -category \mathcal{C} are spaces $\mathcal{C}_{(i)}$ of i -morphisms for $i \leq n$. Basic data of \mathcal{C}^\square are spaces $\mathcal{C}^\square_{(i);(j)}$ for $i, j \leq n$.

Strategy: We will define $\mathcal{C}^\square_{(i);(j)}$ to be the space of diagrams in \mathcal{C} of shape $\Theta^{(i);(j)}$, generalizing $\Theta^{(k);(0)} = \Theta^{(0);(k)} = \Theta^{(k)}$ and $\Theta^{(1);(1)} = \begin{array}{ccc} \bullet & \xrightarrow{\quad} & \bullet \\ \downarrow & \nearrow & \downarrow \\ \bullet & \xrightarrow{\quad} & \bullet \end{array}$ (c.f. §5).

Thm (JF–S): \exists unique sequence of diagrams $\Theta^{(i);(j)}$, the *walking (op)lax i-by-j morphism*, with combinatorics like $\Theta^{(i)} \times \Theta^{(j)}$ but no invertible cells, s.t.:

- There are inclusions $s_h, t_h : \Theta^{(i-1);(j)} \rightrightarrows \Theta^{(i);(j)}$ and $s_v, t_v : \Theta^{(i);(j-1)} \rightrightarrows \Theta^{(i);(j)}$ called “horizontal source and target” and “vertical source and target”.

- The source of the unique $(i + j)$ -dimensional cell $\theta_{i;j}$ is $s_v(\theta_{i;j-1}) \circ s_h(\theta_{i-1;j})$, if i is even, or $t_v(\theta_{i;j-1}) \circ s_h(\theta_{i-1;j})$, if i is odd. (Exactly one of these compositions makes sense.) The target is either $t_h(\theta_{i-1;j}) \circ t_v(\theta_{i;j-1})$ or $t_h(\theta_{i-1;j}) \circ s_v(\theta_{i;j-1})$.

Thm (JF–S): The spaces $\mathcal{C}^\square_{(i);(j)} = \text{maps}(\Theta^{(i);(j)}, \mathcal{C})$ do in fact package into a double (∞, n) -category.

The nontrivial part is to prove completeness.

4. Twisted field theory

Thm (JF–S): If \mathcal{C} is symmetric monoidal, so is \mathcal{C}^\square .

Defn: Choose \mathcal{C} a $\text{sym} \otimes (\infty, n + 1)$ -category, e.g. $\mathcal{C} = \text{VECT}_{n+1}$. Choose $\text{sym} \otimes (\infty, n)$ -category SPACETIMES controlling n -dim field theory.

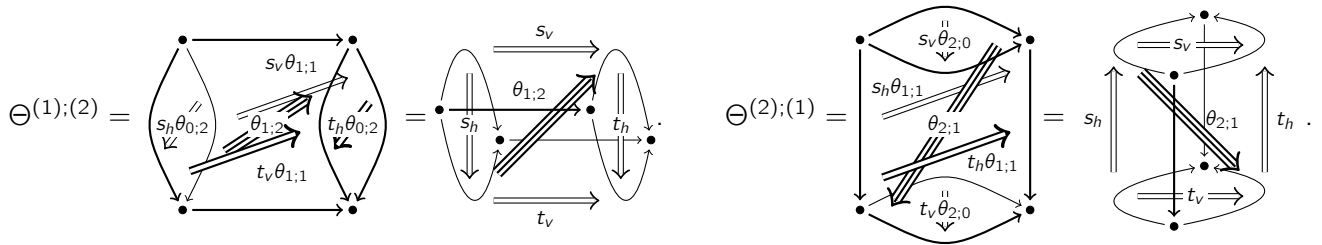
Given $\mathcal{Z} : \text{SPACETIMES} \rightarrow \mathcal{C}$, a *lax \mathcal{Z} -twisted field theory* is a sym monoidal functor $Z : \text{SPACETIMES} \rightarrow \mathcal{C}^\downarrow$ with $s_v \circ Z = \mathbb{1}$ and $t_v \circ Z = \mathcal{Z}$. For *oplax*, use \mathcal{C}^\rightarrow .

Thm (JF–S): Lax $\mathbb{1}$ -twisted field theories are canonically equivalent to “Schrödinger picture” field theories valued in $\Omega\mathcal{C} = \text{sym} \otimes (\infty, n)$ -category of endomorphisms of $\mathbb{1} \in \mathcal{C}$. (If $\mathcal{C} = \text{VECT}_{n+1}$, then $\Omega\mathcal{C} = \text{VECT}_n$.)

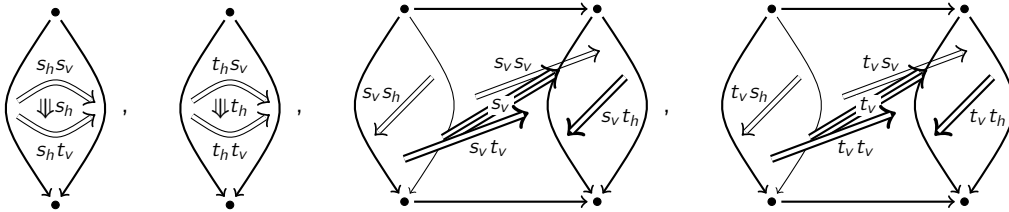
Not true for “lax” \rightsquigarrow “oplax” or “strong.”

Remark: Oplax (but not lax) \mathcal{Z} -twisted field theories are “valued in \mathcal{Z} ” (c.f. §6). Your choice: is “valued in \mathcal{Z} ” or “trivially-twisted = untwisted” more important?

5. Examples of walking morphisms

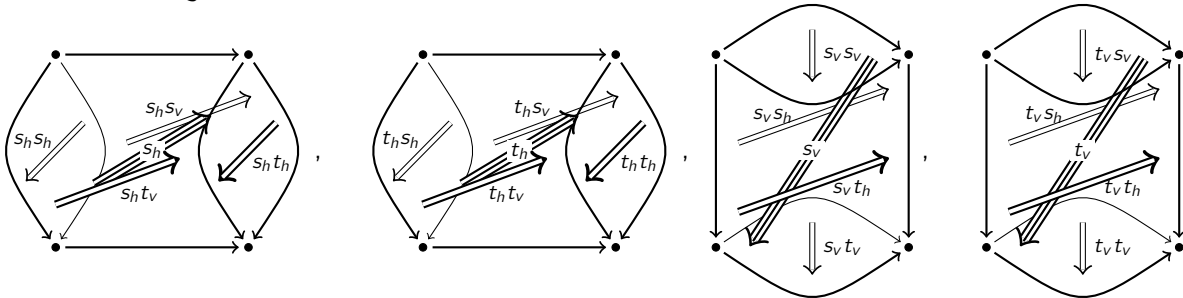


$\Theta^{(1):(3)}$ has the following four faces:

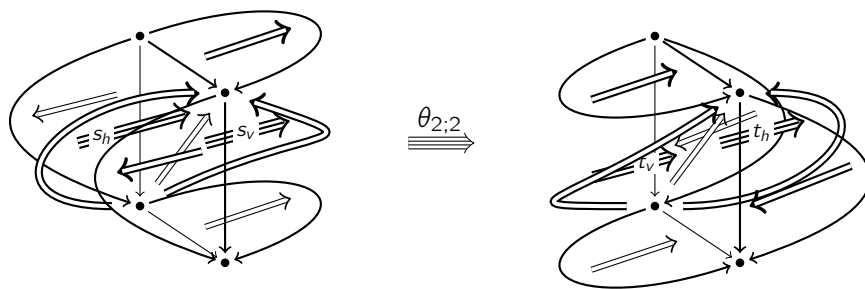


These are glued together along $t_v t_v = s_v t_v$, $t_v s_v = s_v s_v$, $s_v s_h = s_h s_v$, $t_v s_h = s_h t_v$, $s_v t_h = t_h s_v$, and $t_v t_h = t_h t_v$. There is a 4-morphism $\theta_{1,3} : t_v \circ s_h \Rightarrow t_h \circ s_v$.

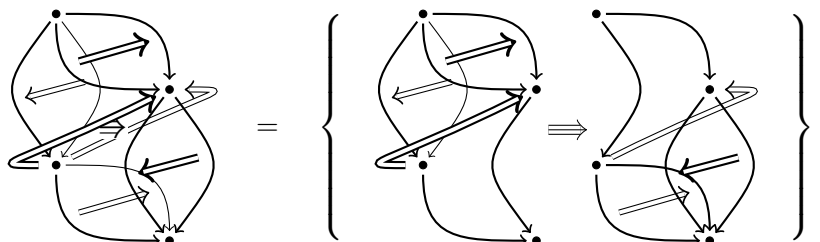
$\Theta^{(2):(2)}$ has the following four faces:



These are glued along $s_h s_h = t_h s_h$, $s_h t_h = t_h t_h$, $s_v s_v = t_v s_v$, $s_v t_v = t_v t_v$, $s_h s_v = s_v s_h$, $s_h t_v = t_v s_h$, $t_h s_v = s_v t_h$, and $t_h t_v = t_v t_h$. Completing $\Theta^{(2):(2)}$ is a 4-morphism $\theta_{2,2} : s_v \circ s_h \Rightarrow t_h \circ t_v$ filling in the 4-ball with boundary this glued-up 3-sphere:

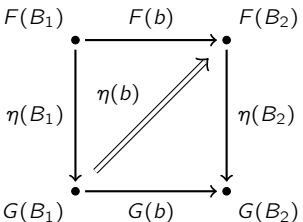
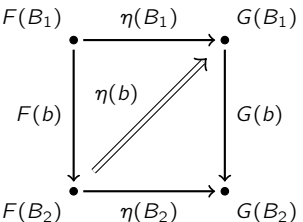
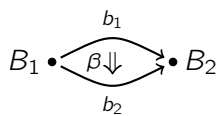
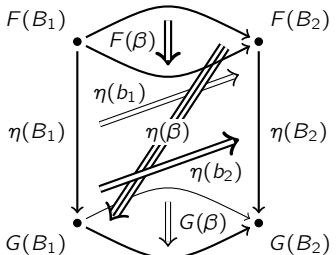
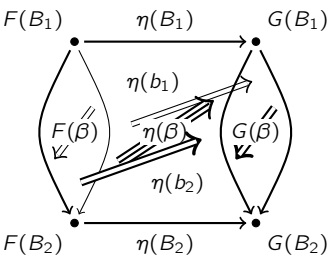


In spite of the distortions, $s_v \circ s_h$ and $t_h \circ t_v$ are parallel 3-morphisms — both are of shape

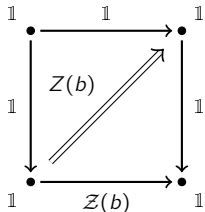
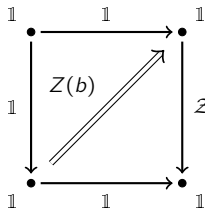


6. Summary: lax versus oplax transformations; lax versus oplax twisted field theory

Lax and oplax transformations $\eta : F \Rightarrow G$ between functors $F, G : \mathcal{B} \Rightarrow \mathcal{C}$ make the following assignments:

		lax	oplax
object	$B \in \mathcal{B}$	$\eta(B) : F(B) \rightarrow G(B)$	$\eta(B) : F(B) \rightarrow G(B)$
1-morphism	$\{B_1 \xrightarrow{b} B_2\} \in \mathcal{B}_{(1)}$		
2-morphism			

Therefore a lax/oplax twisted qft $Z : \mathbb{1} \rightarrow \mathcal{Z}$ assigns to closed manifolds:

		lax	oplax
object	$b \in \text{SPACETIMES}_{(0)}$	$Z(b) : \mathbb{1} \rightarrow \mathcal{Z}(b)$	$Z(b) : \mathbb{1} \rightarrow \mathcal{Z}(b)$
1-morphism	$\mathbb{1} \xrightarrow{b} \mathbb{1} \in \text{SPACETIMES}_{(1)}$		
2-morphism	