

How to Build a Hopf Algebra

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Based on [arXiv:2508.16787](https://arxiv.org/abs/2508.16787), joint with David Reutter

These slides:

<https://categorified.net/Hopf-SimonsFoundation.pdf>



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Statement of the Main Theorem

Main Theorem (JF-Reutter): Let $(\mathcal{C}, 1_{\mathcal{C}} \in \mathcal{C})$ be a pointed $(\infty, 3)$ -category. Given a retract $(1_{\mathcal{C}} \xrightarrow{f} X \xrightarrow{g} 1_{\mathcal{C}}, \alpha : \text{id}_{1_{\mathcal{C}}} \xrightarrow{\cong} gf)$, if f is a left adjoint, g is a right adjoint, and the canonical 2-morphism $\alpha^{\curvearrowright} : f^R \Rightarrow g$ is a right adjoint (equiv. canonical 2-morphism $\alpha^{\curvearrowright} : g^L \Rightarrow f$ is a left adjoint), then

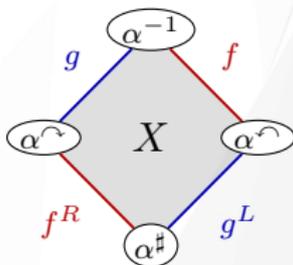
$$\begin{array}{ccccc}
 1_{\mathcal{C}} & \xlongequal{\quad} & 1_{\mathcal{C}} & \xlongequal{\quad} & 1_{\mathcal{C}} \\
 \parallel & & \downarrow & & \parallel \\
 & \nearrow \alpha^{\curvearrowright} & f & \nearrow \alpha^{-1} & \\
 1_{\mathcal{C}} & \xrightarrow{g^L} & X & \xrightarrow{g} & 1_{\mathcal{C}} \\
 \parallel & & \downarrow & & \parallel \\
 & \nearrow ((\alpha^{\curvearrowright})^L)^{\curvearrowright} & f^R & \nearrow \alpha^{\curvearrowright} & \\
 & \parallel \text{R} & & & \\
 & \nearrow ((\alpha^{\curvearrowright})^R)^{\curvearrowright} & & & \\
 1_{\mathcal{C}} & \xlongequal{\quad} & 1_{\mathcal{C}} & \xlongequal{\quad} & 1_{\mathcal{C}}
 \end{array}$$

$$\in \Omega^2 \mathcal{C} := \text{End}_{\text{End}_{\mathcal{C}}(1_{\mathcal{C}})}(\text{id}_{1_{\mathcal{C}}}).$$

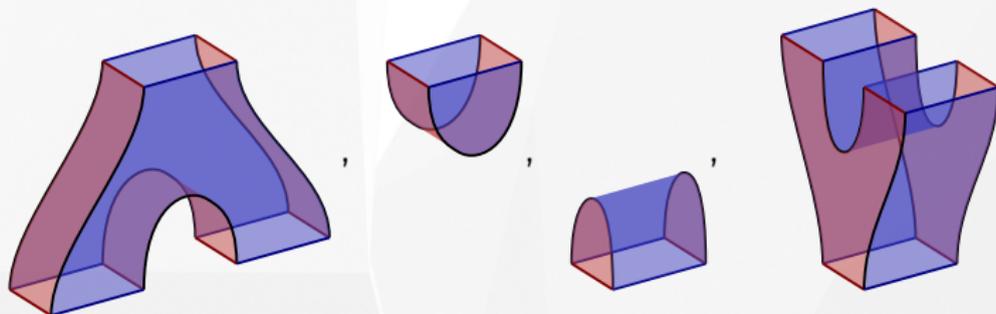
is a **bialgebra**, and its opposite bialgebra is **Hopf** (admits an antipode which is **not necessarily invertible**).

Incoherent topological unpacking of Main Theorem

Poincaré-dually, the composite is a square with alternating boundary conditions:



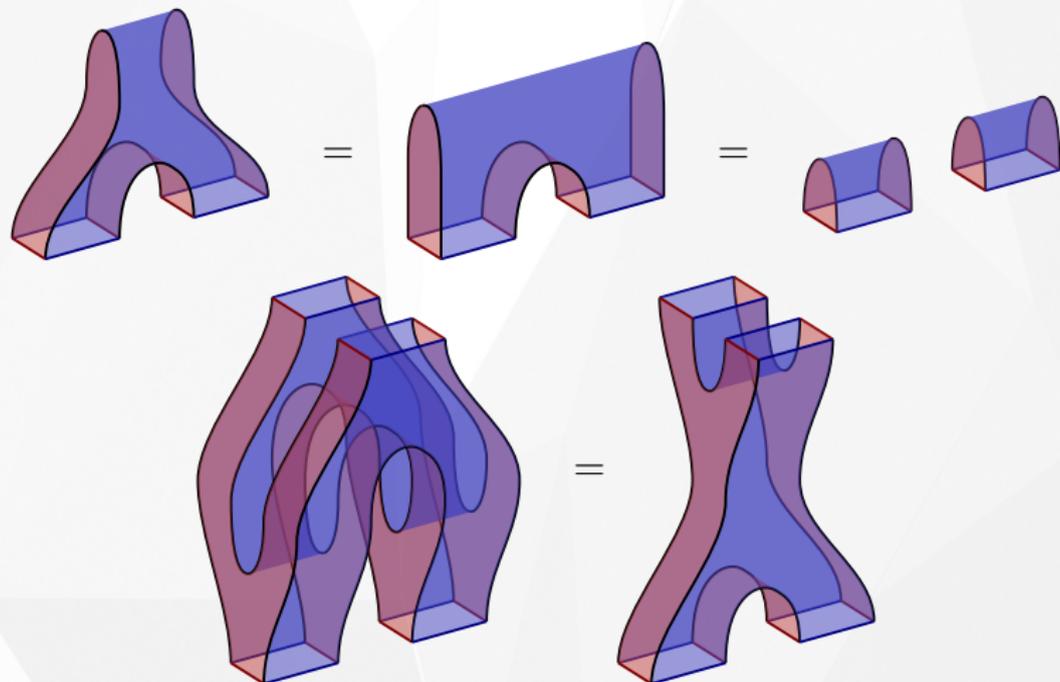
The multiplication, unit, counit, and comultiplication are:



Compare: Reutter 2017, Freed–Teleman 2022, Dimofte–Niu 2024.

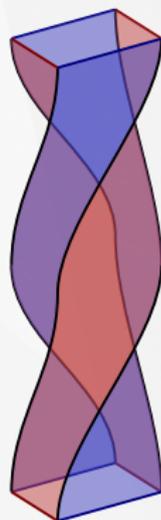
Incoherent topological description of Hopf axioms

Graphically, the retract axiom $\text{id}_{1c} \xrightarrow{\approx} gf$ says that bigons, and bigonal holes, can be created and destroyed.



Dangers of incoherent topological unpacking

- ▶ Tracking framing data is annoying.
- ▶ Calculus of defects not rigorously established.
- ▶ Bialgebra axioms hold up to homotopy for top'l reasons. Want them ∞ -coherently.
- ▶ Want $\mathbf{Ret}^{\text{adj}}(\mathcal{C}) \rightarrow \mathbf{Hopf}(\Omega^2\mathcal{C})$ as a functor.
- ▶ We do not require full adjointibility. Without full adjointibility, cannot produce the antipode as a bordism.
- ▶ This is a feature: we can get every Hopf algebra in every presentably braided monoidal $(\infty, 1)$ -category, including ones with noninvertible antipode.



What the antipode wants to be, but isn't without full adjointibility.

Inductive definition:

$$\mathbf{Cat}_{(\infty, n)} := \{\text{categories enriched in } \mathbf{Cat}_{(\infty, n-1)}\}.$$

$\mathbf{Cat}_{(\infty, \infty)}$ is the $n \rightarrow \infty$ limit. Technically:

$$\mathbf{Cat}_{(\infty, \infty)} := \mathbf{Sheaves} \left(\bigcup_{n \rightarrow \infty} \mathbf{Cat}_{(\infty, n)} \right).$$

Picture (∞, ∞) -categories as **oriented cell complexes**, with a k -cell for each generating k -morphism, oriented from source to target.

The lax tensor product \otimes

Want to orient a product of oriented cell complexes. I.e. given cells $a \in A$, $b \in B$, which direction points $a \otimes b \in A \otimes B$? Strategy: Oriented cell complex $A \rightsquigarrow$ chain complex $C_*(A)$, via $\partial a = \text{target}(a) - \text{source}(a)$; ask $C_*(A \otimes B) = C_*(A) \otimes C_*(B)$, of course with $\partial(a \otimes b) = \partial a \otimes b + (-1)^{\dim a} a \otimes \partial b$.

Theorem (Campion): $\mathbf{Cat}_{(\infty, \infty)}$ admits a unique closed monoidal structure \otimes compatible with this strategy.

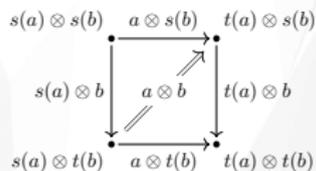
This \otimes is not symmetric: $A \otimes B \not\cong B \otimes A$ in general. Rather, $B \otimes A = (A^{\text{op}} \otimes B^{\text{op}})^{\text{op}}$, where $(-)^{\text{op}}$ takes opposites in all odd directions.

If $A \in \mathbf{Cat}_{(\infty, m)}$ and $B \in \mathbf{Cat}_{(\infty, n)}$, then $A \otimes B \in \mathbf{Cat}_{(\infty, m+n)}$.

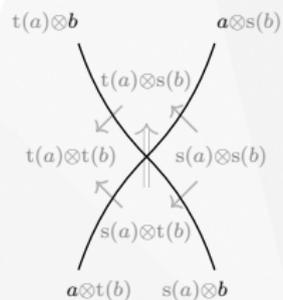
Pictures of \otimes

$$(s(a) \xrightarrow{a} t(a)) \otimes (s(b) \xrightarrow{b} t(b))$$

Pasting diagram



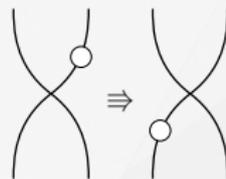
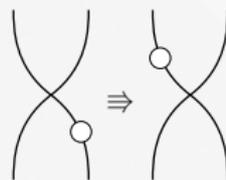
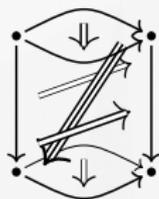
Poincaré dual



arrow \otimes 2-cell



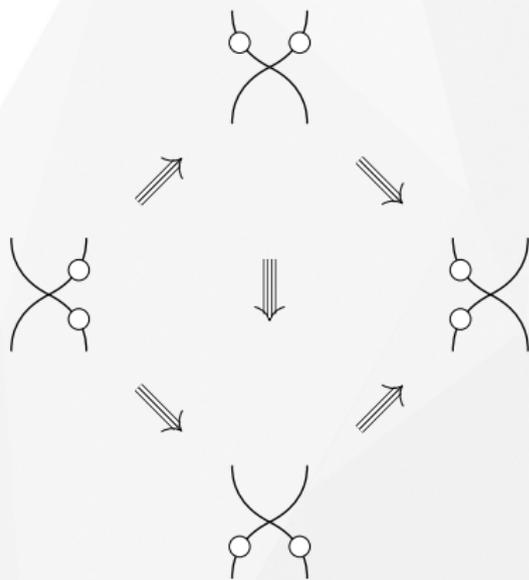
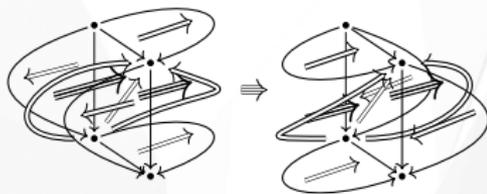
2-cell \otimes arrow



(strictly speaking...)

More pictures of \otimes

2-cell \otimes 2-cell



Lax and oplax natural transformations

If $f, g : A \rightarrow B$ are functors of (∞, ∞) -categories, a **strong natural transformation** $\eta : f \Rightarrow g$ chooses a morphism $\eta a : fa \rightarrow ga$ for each object $a \in A$, and for each $\alpha : a \rightarrow a'$, a naturality isomorphism $(g\alpha) \circ (\eta a) \cong (\eta a') \circ (f\alpha)$. The idea of **(op)lax natural transformations** is to allow noninvertible naturality morphisms. $\text{Fun}^{(\text{op})\text{lax}}(X, Y) :=$ functors $X \rightarrow Y$ and (op)lax natural transfor. These are defined by hom-tensor adjunction:

$$\begin{aligned}\text{maps}(W \rightarrow \text{Fun}^{\text{lax}}(X, Y)) &= \text{maps}(X \otimes W \rightarrow Y) \\ \text{maps}(W \rightarrow \text{Fun}^{\text{oplax}}(X, Y)) &= \text{maps}(W \otimes X \rightarrow Y).\end{aligned}$$

Further ingredients: localization, \otimes , and Ω

$L_n : \mathbf{Cat}_{(\infty, \infty)} \rightarrow \mathbf{Cat}_{(\infty, n)}$, the left adjoint to the inclusion $\mathbf{Cat}_{(\infty, n)} \subset \mathbf{Cat}_{(\infty, \infty)}$, inverts all ($> n$)-morphisms.

$\mathbf{Cat}_{(\infty, \infty)}^* := \{\text{categories equipped with a distinguished object}\}$

The **smash product** of $(X \ni 1_X), (Y \ni 1_Y) \in \mathbf{Cat}_{(\infty, \infty)}^*$ is

$$X \otimes Y := \frac{X \otimes Y}{X \otimes \{1_Y\} \sqcup_{\{1_X\} \otimes \{1_Y\}} \{1_X\} \otimes Y}$$

\otimes is hom-tensor adjoint to **strongly pointed** functors — $f : X \rightarrow Y$ with $f1_X \cong 1_Y$ — and their (op)lax natural transformations.

$\vec{S}^1 := \mathbf{BN}$ is the **directed circle**. The **based loop category** of $X \in \mathbf{Cat}_{(\infty, \infty)}^*$ is $\Omega X := \mathbf{Fun}_*^{\text{oplax}}(\vec{S}^1, X) = \mathbf{End}_X(\text{basepoint})$.

Monad \otimes Monad = Bialgebra

A **braided monoidal** $(\infty, 1)$ -category is a monoid object in $\{\text{monoidal } (\infty, 1)\text{-categories}\}$. If $\mathcal{C} \in \mathbf{Cat}_{(\infty, 3)}^*$, then $\Omega\mathcal{C}$ is a mon. $(\infty, 2)$ -cat. and $\Omega^2\mathcal{C}$ is a braided mon. $(\infty, 1)$ -cat.

Vague idea: **(op)lax monoidal functor** is some type of **(co)algebra**.

$\mathbf{Mnd} := \mathbf{B}(\text{free monoidal 1-category on an associative algebra})$. In other words, $\text{maps}_*(\mathbf{Mnd}, X) = \text{ob Alg}(\Omega X)$. Define **oplax algebra morphisms** by $\text{Alg}^{\text{oplax}}(\Omega X) := \text{Fun}_*^{\text{oplax}}(\mathbf{Mnd}, X)$.

Theorem (Hadzihasanovic, JF-R): If $X \in \mathbf{Cat}_{(\infty, 3)}^*$, then $\text{Alg}^{\text{oplax}}(\text{Alg}^{\text{oplax}}(\Omega X)) = \text{BiAlg}(\Omega^2 X)$. Equivalently:

$$L_3(\mathbf{Mnd} \otimes \mathbf{Mnd}) = \mathbf{B}^2(\text{free braided 1-category on a bialgebra}).$$

Incoherent pictures of $\text{Mnd} \otimes \text{Mnd} = \text{Bialg}$

Can think of $\text{Arr} \otimes \text{Arr}$ as a pasting square. Poincaré dually, it is two wires crossing. Writing in grey the cells that are smashed out, here is a Poincaré dual picture of the generating cell in $\vec{S}^1 \otimes \vec{S}^1$:

$$B = \begin{array}{c} \diagup \quad \diagdown \\ \bullet \\ \diagdown \quad \diagup \end{array}$$

Mnd is generated by \vec{S}^1 and a multiplication 2-cell m (and a unit, omitted). Multiplying m with the generating 1-cell in Mnd gives the two generating 3-cells in $\text{Mnd} \otimes \text{Mnd}$:

$$\underbrace{\begin{array}{c} \text{Diagram 1} \\ \text{Diagram 2} \end{array}}_{m_B: B \otimes B \rightarrow B} \quad \text{and} \quad \underbrace{\begin{array}{c} \text{Diagram 3} \\ \text{Diagram 4} \end{array}}_{\Delta: B \rightarrow B \otimes B}$$

The bialgebra law $\Delta \circ m = (m \otimes m) \circ (\text{permutation}) \circ (\Delta \otimes \Delta)$ comes from the 4-cell $m \otimes m \in \text{Mnd} \otimes \text{Mnd}$.

(-) \otimes Adjunction and Adjunction \otimes (-)

Adj := free 2-category on an adjunction. There is a canonical inclusion $\text{Mnd} \hookrightarrow \text{Adj}$ that restricts to endomorphisms of the source of the left adjoint. There is a canonical epimorphism $\ell : \text{Arr} \twoheadrightarrow \text{Adj}$ that imposes that the arrow is a left adjoint.

Theorem (JF–Scheimbauer, JF–R, Masuda):

- ▶ $\text{id} \otimes \ell : k\text{-Cell} \otimes \text{Arr} \rightarrow k\text{-Cell} \otimes \text{Adj}$ is the epimorphism that imposes that $\partial(k\text{-Cell}) \otimes \text{Arr}$ factors through $\partial(k\text{-Cell}) \otimes \text{Adj}$ and also that the $(k + 1)$ -dimensional filling is a left adjoint.
- ▶ $\ell \otimes \text{id} : \text{Arr} \otimes k\text{-Cell} \rightarrow \text{Adj} \otimes k\text{-Cell}$ is the epimorphism that imposes that $\text{Arr} \otimes \partial(k\text{-Cell}) \rightsquigarrow \text{Adj} \otimes \partial(k\text{-Cell})$ and also that the **mate** of the $(k + 1)$ -dim filling is a left adjoint.

Provided f, g are left adjoints, the **mate** of $\phi : f \Rightarrow g$ is the canonical map $\phi^{\curvearrowright} : g^R \Rightarrow f^R$.

Adj \otimes Adj \rightsquigarrow Retract \rightsquigarrow coherence proof

Corollary: $L_3(\text{Adj} \otimes \text{Adj}) =$ the free 3-category on a retract $(1_c \xrightarrow{f} X \xrightarrow{g} 1_c, \alpha : \text{id}_{1_c} \xrightarrow{\sim} gf)$ in which

- ▶ f is a left adjoint and g is a right adjoint,
- ▶ and the canonical 2-morphism $\alpha^{\curvearrowright} : f^R \Rightarrow g$ is a right adjoint, equiv. canonical 2-morphism $\alpha^{\curvearrowleft} : g^L \Rightarrow f$ is a left adjoint.

Main coherence statement: $\text{Mnd} \hookrightarrow \text{Adj}$ and so

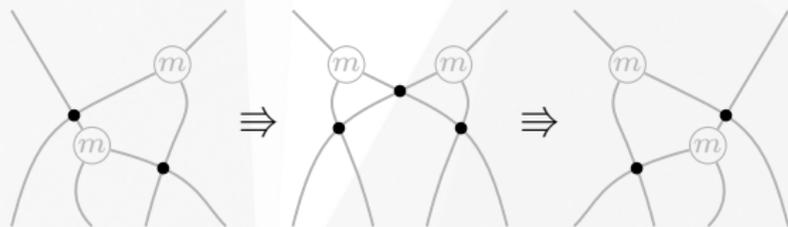
$$\begin{array}{ccc} \overbrace{L_3(\text{Mnd} \otimes \text{Mnd})}^{\text{walking bialgebra}} & \rightarrow & \overbrace{L_3(\text{Adj} \otimes \text{Adj})}^{\text{walking adjunctible retract}} \\ \underbrace{\text{Fun}^{\text{oplax}}(L_3(\text{Mnd} \otimes \text{Mnd}), \mathcal{C})}_{\text{BiAlg}(\Omega^2 \mathcal{C})} & \leftarrow & \underbrace{\text{Fun}^{\text{oplax}}(L_3(\text{Adj} \otimes \text{Adj}), \mathcal{C})}_{\text{Ret}^{\text{adj}}(\mathcal{C})} \end{array}$$

The antipode axiom

A bialgebra B is **Hopf** if and only if the **shear map**

$$B \otimes B \xrightarrow{\Delta_B \otimes \text{id}_B} B \otimes B \otimes B \xrightarrow{\text{id}_B \otimes m_B} B \otimes B$$

is invertible. In $\text{Mnd} \triangleleft \text{Mnd}$, the **universal** shear map looks like:

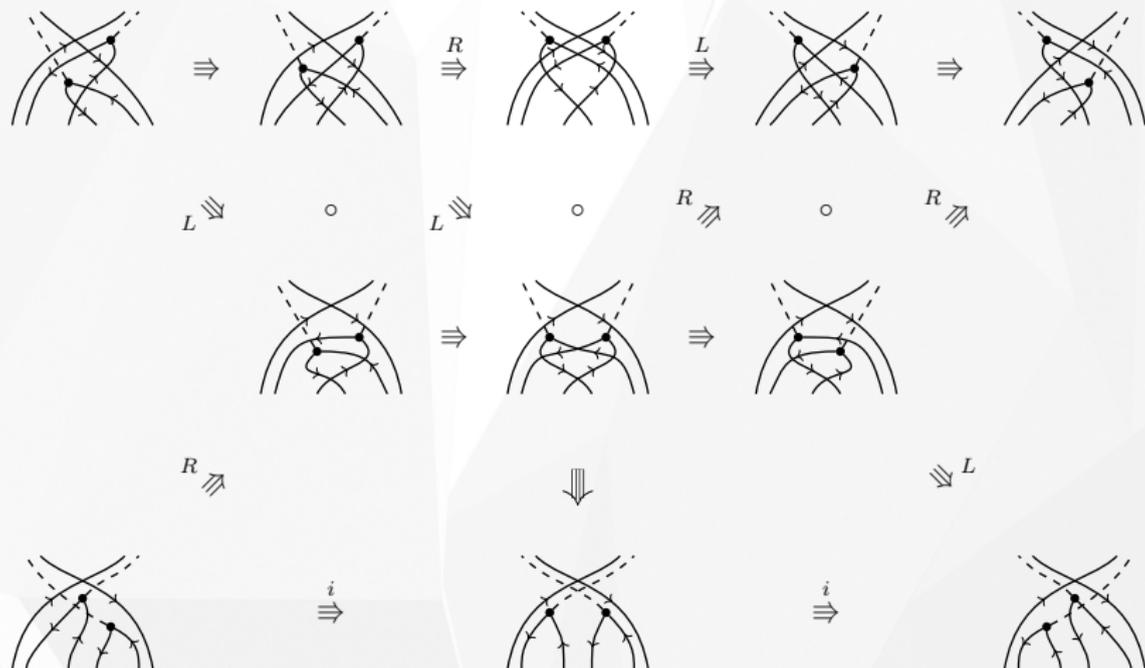


The bialgebra B is the intersection point. The light gray shows the cells in $\text{Mnd} \otimes \text{Mnd}$ that are smashed out in $\text{Mnd} \triangleleft \text{Mnd}$.

Theorem: The universal shear map inverts when sent via $\text{Mnd} \otimes \text{Mnd} \rightarrow \text{Adj} \otimes \text{Adj}$ and 3-localized.

Proof of existence of antipode

Theorem: The universal shear map inverts when sent via $\text{Mnd} \otimes \text{Mnd} \rightarrow \text{Adj} \otimes \text{Adj}$ and 3-localized. **Proof:**



(plus some technical manoeuvres with orientals in order to justify the use of strict graphical calculus)